

Models of linear dependent type theory

Martin Lundfall

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Abstract

In this paper, we construct a type theory which deals with non-linear, ordinary dependent types (which we will call *cartesian*), and *linear types*, where both constructs may depend on terms of cartesian types. In the interplay between the cartesian and linear types we introduce the new type formers $\prod_{x:A} B$ and $\prod_{x:A} B$, akin to Π and Σ , but where the dependent type B , (and therefore the resulting construct) is a linear type. These can be seen as internalizing quantification over linear propositions. We also introduce the modalities M and L , transforming linear types into cartesian types and vice versa.

We interpret the theory in a split comprehension category $\pi : \mathcal{T} \rightarrow \mathcal{C}^{\rightarrow}$ [Jac93], accompanied by a split monoidal fibration (Definition 4.3), $q : \mathcal{L} \rightarrow \mathcal{C}$. The intuition is that \mathcal{C} models a category of contexts, so that for any $\Gamma \in \mathcal{C}$, the fiber \mathcal{T}_{Γ} is the category containing the cartesian types which can be formed in the context Γ , while the fiber \mathcal{L}_{Γ} is a monoidal category of linear types in Γ . In this setting, the type formers $\prod_{x:A}$ and $\prod_{x:A}$ are understood as right and left adjoints of the monoidal reindexing functor $\pi_A^* : \mathcal{L}_{\Gamma} \rightarrow \mathcal{L}_{\Gamma.A}$ corresponding to the weakening projection $\pi_A : \Gamma.A \rightarrow \Gamma$ in \mathcal{C} . The operators M and L give rise to a fiberwise adjunction $L \dashv M$ between \mathcal{L} and \mathcal{T} , where we understand the traditional exponential modality as the comonad $! = LM$.

We provide two concrete examples of models, the *set-indexed families model* and the *diagrams model*. In the former, cartesian types are interpreted in the familiar way, as sets indexed by their context set Γ , and linear types are interpreted as Γ -indexed family of objects of a symmetric monoidal category \mathcal{V} . The latter model extends the groupoid model of dependent type theory [HS98] to accommodate linear types. Here, cartesian types over a context Γ are interpreted as a family of groupoids indexed over the groupoid Γ , while linear types are interpreted as diagrams over groupoids, $A : \Gamma \rightarrow \mathcal{V}$ in any symmetric monoidal category \mathcal{V} . We show that the *diagrams model* can under certain conditions model a linear analogue of the univalence axiom, and provide some discussion on the higher-dimensional nature of linear dependent types.