Tentamensskrivning i Mathematics of Cryptography, 7,5 hp 18 March 2019 9.00-14.00

There are ten problems, each giving between 0 and 8 points. The points from the exam is added to points from the homework assignments. Grades are then given by the following intervals:

A 100-92p, B 91-84p, C 83-76p, D 75-68p, E 67-60p.

Remember to motivate your answers carefully. No calculators or computers may be used.

- 1. a) Explain what a public key cryptosystem is. In particular, explain what a one-way-function with a trapdoor is, and its function in a public key cryptosystem.
 - b) Describe the problem one needs to solve to decrypt a ciphertext of the cryptosystem RSA (without knowing the private key). What is the underlying hard mathematical problem for RSA?
 - c) Explain in some detail which basic properties one wants from a public key cryptosystem? 4 p
- 2. a) Let a, b, m be positive integers. Give a criterion in terms of a, b, m for there to be a solution to the equation

$$ax \equiv b \mod m.$$

b) Take
$$a, b, m$$
 so that there is at least one solution to the equation,

$$ax \equiv b \mod m.$$

Give an expression (in terms of a, b, m) for how many solutions there are modulo m. 3 p

c) Use the Chinese reminder theorem to solve the following system of equations:

$$\begin{cases} 7x \equiv 3 \mod 10\\ 7x \equiv 8 \mod 27 \end{cases}$$

 $2\,\mathrm{p}$

 $5\,\mathrm{p}$

 $1\,\mathrm{p}$

- 3. a) Let G be a finite group and g an element of G. Prove how many multiplications that are needed, using fast powering, to compute g^N in terms of the positive integer N. Is the growth polynomial/subexponential/exponential with the input?
 - b) Say that p = 23, q = 56499605716734596849, n = pq and that a is a primitive root both modulo p and modulo q. Roughly how many steps would it take for Pollard's p-1-algorithm, using the integer a as a base, to give the factorization of n? 3 p
- 4. a) State the ElGamal problem and state the Diffie-Hellman problem. 2 p
 - b) Is the ElGamal problem easier to solve than the Diffie-Hellman problem? Is it harder? (No proof is necessary.)
 - c) What are the algorithms involved in encryption and decryption in the ElGamal cryptosystem? What are their complexity? Are they polynomial/subexponential/exponential?3 p

 $3\,\mathrm{p}$

 $2\,\mathrm{p}$

 $2\,\mathrm{p}$

| | d) | What is the fastest algorithm (that we know) that breaks the ElGamal cryptosystem? What is its complexity? Is it polynomial/subexponential/exponential? | $2\mathrm{p}$ |
|----|----|--|---------------|
| 5. | a) | Explain what a digital signature scheme is. | $2\mathrm{p}$ |
| | b) | What is the problem that digital signatures are supposed to solve? | $2\mathrm{p}$ |
| | c) | Explain the man-in-the-middle attack against a public key cryptosystem (of your choice). | $2\mathrm{p}$ |
| | d) | Does a digital signature protect against a man-in-the-middle attack? | $2\mathrm{p}$ |
| 6. | a) | Describe the Fermat primality test (that is, a primality test based upon Fermat's little theorem). | $3\mathrm{p}$ |
| | b) | What are the benefits of using the Miller-Rabin primality test compared to the Fermat primality test? | $2\mathrm{p}$ |
| | c) | What are the benefits of using the Fermat primality test compared to the Miller-Rabin primality test? | 1 p |
| | d) | What is the complexity of the Miller-Rabin primality test if one want to use it to get a primality proof? Is is polynomial/subexponential/exponential? | 1 p |
| | e) | Name an algorithm that gives a primality proof that has substantially better complexity than the Miller-Rabin primality test when used to give a primality proof. | 1 p |
| 7. | a) | What is a <i>B</i> -smooth number? | 1 p |
| | b) | Let p, q be prime numbers, $n = pq$, $a = \lfloor \sqrt{n} \rfloor + 1$, and $F(T) = T^2 - n$. Say that we have computed the list of integers $F(a), F(a+1), \ldots, F(a+b)$ for some positive integer b . Describe how the quadratic sieve gives all B -smooth numbers in this list (for any choice of positive integer B). | 4 p |
| | c) | Give an approximate expression for the number of divisions of integers one needs to do in the process described in b). | 1 p |

- d) Let n be of size 2^k . Give an expression for B that grows subexponentially with k, such that the expected number of checks of random integers of size roughly \sqrt{n} one needs to do in order to find $\pi(B)$ integers that are B-smooth also grows subexponentially. What is this expected number?
- 8. a) Consider the elliptic curve over \mathbb{F}_5 given by,

$$E: y^2 = x^3 + 4x + 4.$$

List the points of $E(\mathbb{F}_5)$.

b) Lenstra's factorization algorithm is subexponential. Explain in some detail how this fact depends upon the distribution of B-smooth numbers. $6\,\mathrm{p}$

 $2\,\mathrm{p}$

 $2\,\mathrm{p}$

9. Use index calculus to solve the DLP: $g^x \equiv_p h$ with g = 103, h = 386 and p = 1019. The fact that $hg^{183} = 126$ and the following table will be helpful:

$$\begin{pmatrix} i & g^i \pmod{p} \\ 946 & 2 \cdot 3 \cdot 5 \cdot 7 \\ 735 & 2 \cdot 3^2 \cdot 5 \\ 347 & 2^3 \cdot 3 \\ 245 & 3 \cdot 7 \\ 454 & 2 \cdot 3^2 \cdot 5 \cdot 7 \end{pmatrix}$$

 $8\,\mathrm{p}$

- 10. a) What the expected number of steps for the Pollard- ρ method to find the solution to a DLP in \mathbb{F}_{p}^{*} ? Is this algorithm polynomial/subexponential/exponential? 2 p
 - b) What is the main advantage of the Pollard- ρ method over Shank's Babystep-Giantstep method to solve a DLP? 1 p
 - c) The table

$$\left(\begin{array}{cccccc}
i & x_i & y_i & f(y_i) \\
0 & 1 & 1 & 11 \\
1 & 11 & 6 & 20 \\
2 & 6 & 14 & 12 \\
3 & 20 & 6 & 20 \\
4 & 14 & 14 & \end{array}\right)$$

describes an application of the function

$$f(x) = \begin{cases} gx \pmod{p} & \text{if } 0 \le x < 8\\ x^2 \pmod{p} & \text{if } 8 \le x < 16\\ hx \pmod{p} & \text{if } 16 \le x < 23 \end{cases}$$

where p = 23, g = 11, h = 3, $x_{i+1} = f(x_i)$ and $y_{i+1} = f(f(y_i))$. Use this data to solve the DLP: $g^x \equiv_p h$. 5 p

The exam will be returned 11.00 on Friday the 5th of April in room 410 in house 6. After that it can be collected in room 204 in house 6.