

This exam consists of 6 questions, worth a total of 40 points. Not all questions are equally difficult. You may submit your answers in either English or Swedish. Write clearly and motivate your answers carefully.

Good luck! — Lycka till!

1 Calculate the cardinalities of the following sets, and order them according to cardinality (some may be equal):

- $S_1 := \{f : \mathbb{N} \rightarrow \mathbb{N} \mid f \text{ increasing}\}$
- $S_2 := \{f : \mathbb{N} \rightarrow \mathbb{N} \mid f \text{ decreasing}\}$
- $S_3 := \mathbb{N}^{\mathbb{R}}$
- $S_4 := \{f : \mathbb{R} \rightarrow \mathbb{N} \mid f \text{ increasing}\}$

(Here “increasing” means “non-strictly increasing”, i.e. $f(x) \leq f(y)$ whenever $x \leq y$, and similarly “decreasing” means “non-strictly decreasing”.)

2 Work in ZF. The *order extension principle* is the statement “for every partially ordered set (X, \leq) , there is some total order \leq' on X with $(\leq) \subseteq (\leq')$,” or more concisely “every partial order can be extended to a total order”.

- (a) Show that AC implies the order extension principle. *(You may use AC either directly, or via any of its consequences considered in the course.)*
- (b) Show that the order extension principle implies the restricted form of AC, “for every family of non-empty *finite* sets, there exists some choice function”.

3 Cardinalities of models.

- (a) Show that any consistent theory T in a language L has some model of size $\leq \max(\|T\|, \|L\|)$.
- (b) Give an example of a consistent theory T in a language L with no model of size $< \max(\|T\|, \|L\|)$.
- (c) Must a consistent theory T in a language L have some model of size exactly $\max(\|T\|, \|L\|)$? Prove or give a counterexample.
- (d) Must a consistent theory T in a language L have some model of size $\leq \|T\|$, assuming $\|T\|$ is infinite? Prove or give a counterexample.

4 Show that the following functions are recursive:

- (a) “truncated predecessor” $p : \mathbb{N} \rightarrow \mathbb{N}$, given by $p(0) := 0$, $p(n) := n - 1$ for $n > 0$.
- (b) “truncated subtraction”, aka “monus” $\dot{-} : \mathbb{N}^2 \rightarrow \mathbb{N}$, $m \dot{-} n := \max(m - n, 0)$.

(Recall that the partial recursive functions are generated by projections, the constant 0, successor, composition, primitive recursive definitions, and the (unbounded) minimisation operator.)

5 ZFC is a theory in the countable language $\langle \in \rangle$. So by the Löwenheim–Skolem theorems, if ZFC is consistent, it has some countable model \mathcal{M} . But ZFC proves “there exists an uncountable set”, so this statement holds in \mathcal{M} . Why doesn’t this give a contradiction?

6 The goal of this problem is to show there are uncountably many countable non-standard models of PA, up to isomorphism. Work over the language of arithmetic $L_A := \langle 0, 1, S, +, \times \rangle$. For $n \in \mathbb{N}$, write \bar{n} for the term $S^n(0)$ in L_A . In L_A , define “ $x \geq y$ ” as “ $\exists z x = y + z$ ” and “ $x \mid y$ ” as “ $\exists z y = z \times x$ ”. A model of PA is *standard* if every element is the interpretation of some \bar{n} ; equivalently, if it is isomorphic to the model $\langle \mathbb{N}; 0, 1, S, +, \times \rangle$.

- (a) Consider L_A plus an extra constant symbol c ; let T be the theory consisting of PA, together with the axiom “ $c \geq \bar{n}$ ” for each $n \in \mathbb{N}$. Show that T is consistent; deduce that there exists some non-standard model of PA.
- (b) Write $\mathbb{P} \subseteq \mathbb{N}$ for the set of prime numbers. Show that for every set of primes $X \subseteq \mathbb{P}$, there is some countable model \mathcal{M} of PA with an element $a \in \mathcal{M}$ whose standard prime divisors are precisely X ; that is, such that for each $p \in \mathbb{P}$, $\mathcal{M} \models \bar{p} \mid a$ if and only if $p \in X$.
- (c) Deduce that there are uncountably many isomorphism classes of countable models of PA. (Hint: note that for each such model, only countably many sets of primes can appear as sets of standard prime divisors of some element.)

———— End of exam — Slut på provet —————