

# Computational Biology

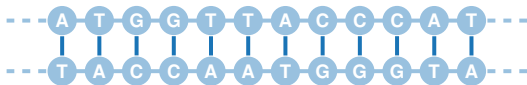
## Ribonucleic Acid (RNA)

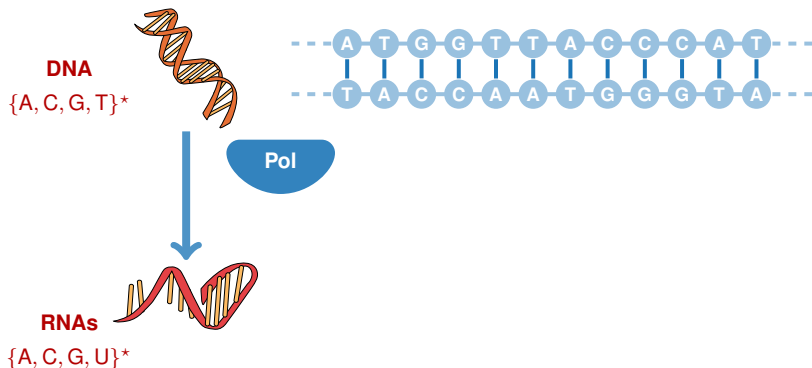
Marc Hellmuth

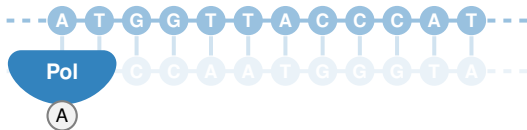
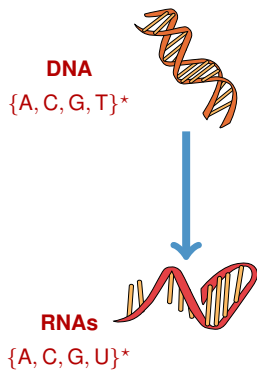
Department of Mathematics  
Stockholm University

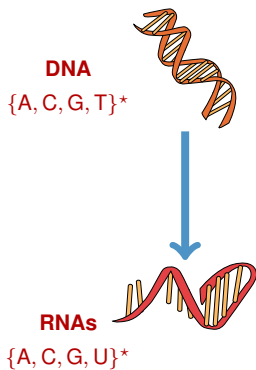
*With kind permission of Yann Ponty and Sebastian Will (l'Ecole Polytechnique LIX), I reuse for the RNA-part many of their slides of the AMI2B course.*

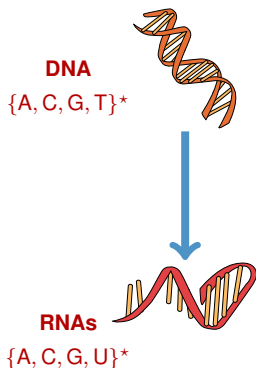
**DNA**  
{A, C, G, T}<sup>\*</sup>

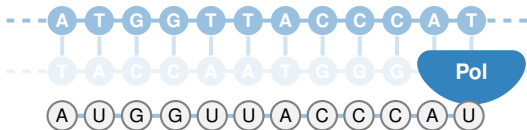
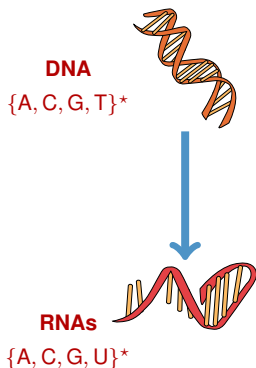


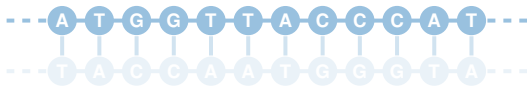
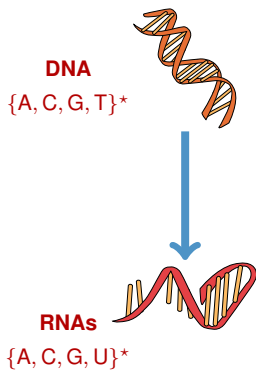




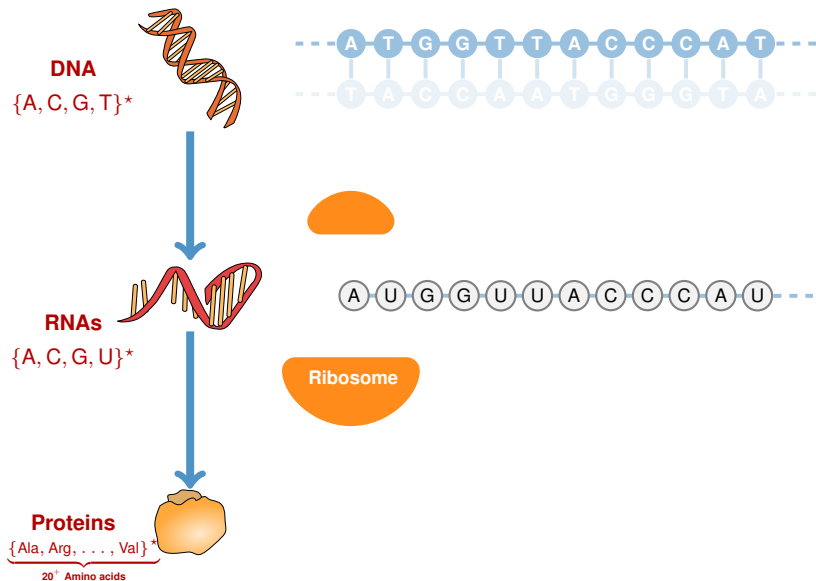


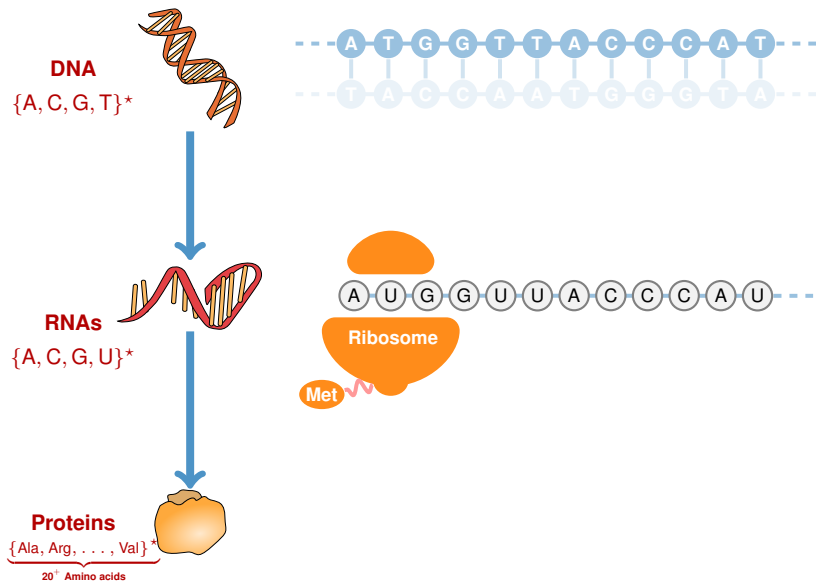


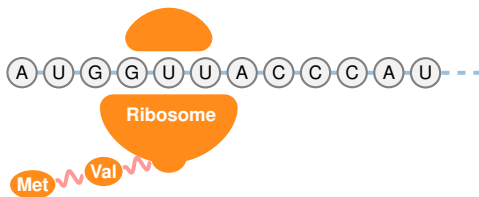
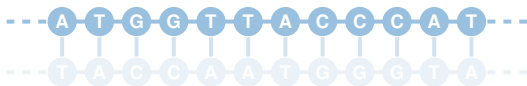
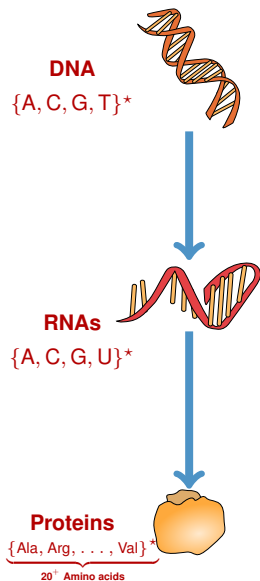


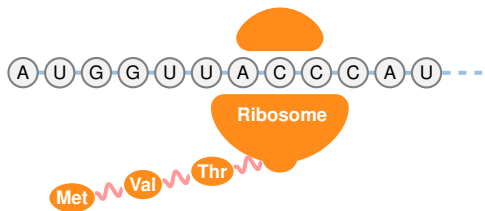
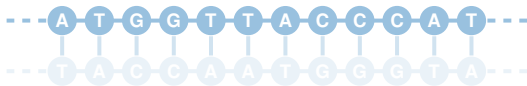
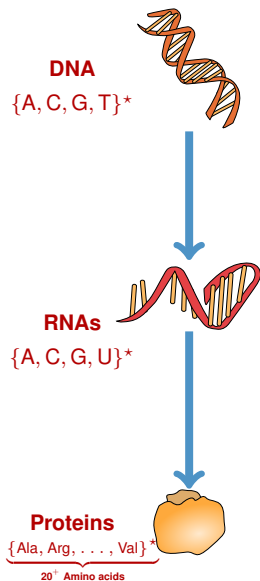


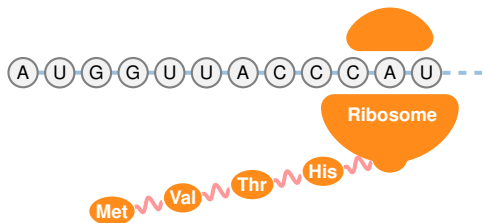
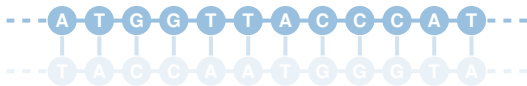
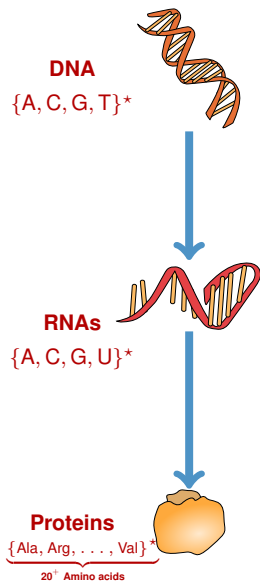


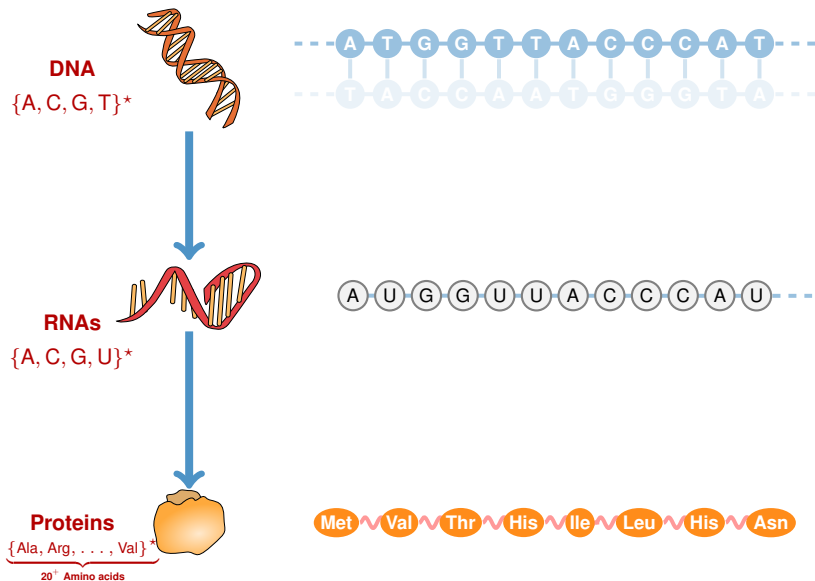


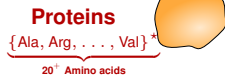
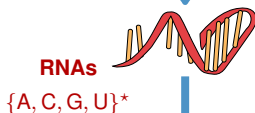
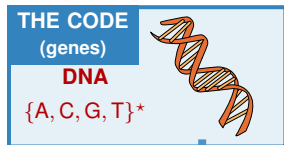


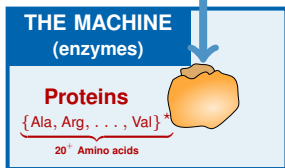
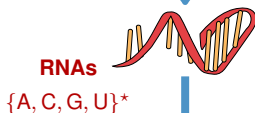
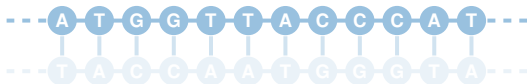
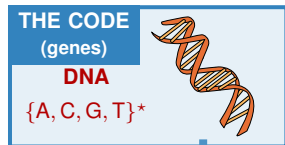















**THE CODE**  
(genes)



**DNA**  
{A, C, G, T}<sup>\*</sup>




**MEH...**



**RNAs**  
{A, C, G, U}<sup>\*</sup>

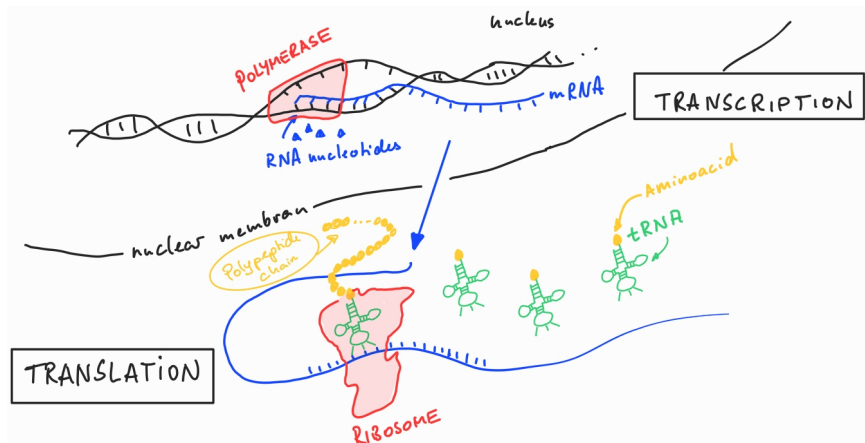


**THE MACHINE**  
(enzymes)



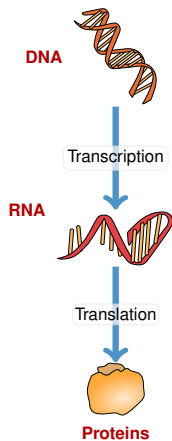
**Proteins**  
{Ala, Arg, . . . , Val}<sup>\*</sup>  
20<sup>+</sup> Amino acids

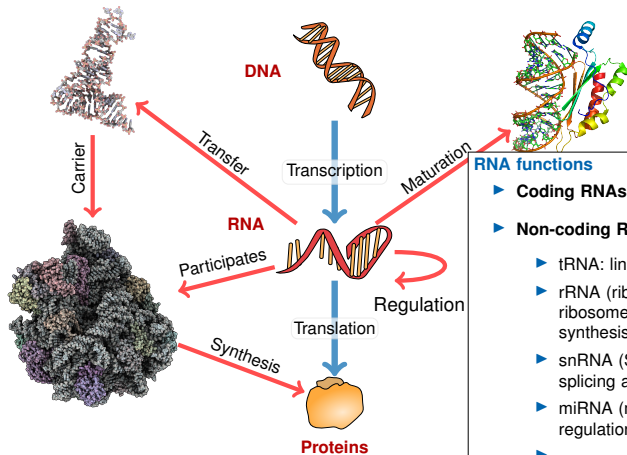




Oh! There are RNAs that are not only messengers ...

great visualization: <https://www.youtube.com/watch?v=gG7uCskUOrA>

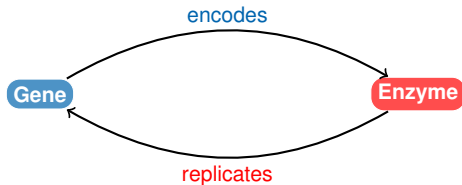




## RNA functions

- ▶ **Coding RNAs** (Translation - mRNA)
- ▶ **Non-coding RNAs**
  - ▶ tRNA: linking codons to aminoacids
  - ▶ rRNA (ribosomalRNA): part of the ribosome, essential for protein synthesis
  - ▶ snRNA (Small nuclear RNA, ~ 150 nt): splicing and other functions
  - ▶ miRNA (microRNA, 21-22 nt): regulation of gene expression
  - ▶ ...

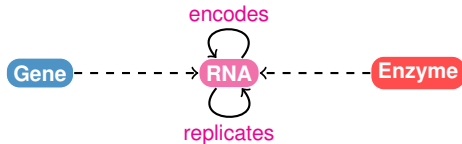
~ 97% are non-coding RNAs in eukaryotes  
 RNA can even act as genome (virus)



A **gene** big enough to specify **an enzyme** would be too big to replicate accurately without the aid of **an enzyme** of the very kind that it is trying to specify. So the system *apparently cannot get started*.

#### RNA - WORLD - HYPOTHESIS:

self-replicating RNA molecules are precursors to all current life on earth.

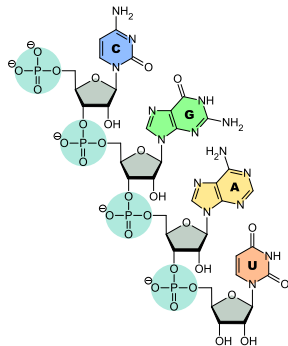


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### RNA - WORLD - HYPOTHESIS:

self-replicating RNA molecules are precursors to all current life on earth.

- ▶ single-stranded polymer
- ▶ polymer made of **nucleotides+backbone**
- ▶ **nucleotides**: guanine (G), adenine (A), uracil (U), cytosine (C)
- ▶ **backbone**: alternating sugar (ribose) and phosphat groups (related to phosphoric acid) nucleotides are attached to sugar
- ▶ the nucleotides of polymer can bind (A-U, C-G, G-U) via hydrogen bonds, i.e., unlike DNA it is more often found in nature as a single-strand folded unto itself, rather than a paired double-strand.

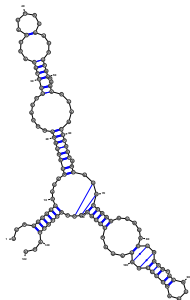


**RNA can fold into complex 3D structures that are essential to its function(s).**

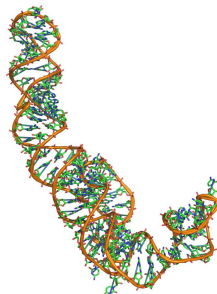
Three\* levels of representation:

```
UUAGGCGGCCACAGC
GGUGGGGUUGCCUCC
CGUACCCAUCCGAA
CACGGAAGUAAGCC
CACCAGCGUCCGGG
GAGUACUGGAGUGCG
CGAGCCUCUGGAAA
CCCGGUUCGCCCA
CC
```

Primary structure



Secondary structure



Tertiary structure

Source: 5s rRNA (PDB 1K73:B)

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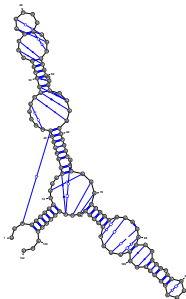
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GAGUACUGGAGUGCG
CGAGCCUCUGGAAA
CCCGUUCGCCGCA
CC
```

Primary structure



Secondary<sup>+</sup> structure



Tertiary structure

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---

\*Well, mostly...

$\mathbb{A} := \{A, C, G, U\}$  and  $\mathbb{B} := \{AU, UA, GC, CG, GU, UG\}$

A **primary structure** (of length  $n$ ) is a sequence  $s = s_1 \dots s_n \in \mathbb{A}^n$ .

A **secondary structure**  $\mathcal{S}$  is a collection of ordered pairs  $(i, j)$ , where  $1 \leq i < j \leq n$ , s.t. the following properties hold:

1. If  $(i, j), (k, l) \in \mathcal{S}$ , then it is not the case that  $i < k < j < l$ .
2. If  $(i, j), (k, l) \in \mathcal{S}$  and  $i \in (k, l)$  implies that  $i = k$  and  $j = l$ .
3. If  $(i, j) \in \mathcal{S}$ , then  $j > i + \theta$ , where  $\theta$  is a fixed integer and usually taken to be 3.

A secondary structure  $\mathcal{S}$  for a given sequence  $s = s_1 \dots s_n \in \mathbb{A}^n$  is a secondary structure fulfilling in addition

4. If  $(i, j) \in \mathcal{S}$ , then  $s_i s_j \in \mathbb{B}$ .

If item 4. is fulfilled, then we say that the sequence  $s \in \mathbb{A}^n$  **realizes**  $\mathcal{S}$ .

A **tertiary structure** is basically the 3D structure, i.e., refers to locations of the atoms in three-dimensional space.

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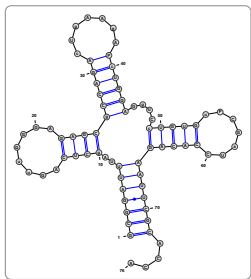
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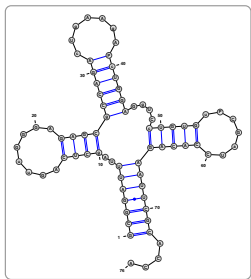


Outer-planar graphs  
Hamiltonian-path,  
 $\Delta(G) \leq 3$ , 2-connected\*

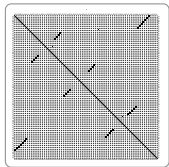
### Supporting intuitions

Different representations  
Common combinatorial structure

\* Additional steric constraints



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Hamiltonian-path,  
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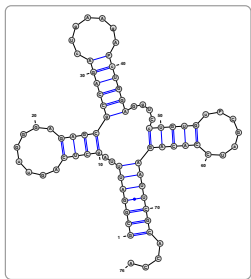


Dot plots  
Adjacency matrices\*

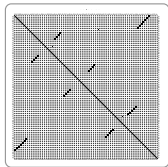
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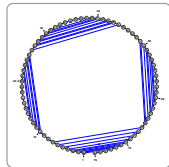
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Outer-planar graphs  
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Dot plots  
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Non-crossing arc diagrams\*

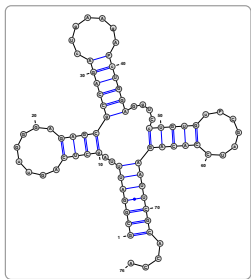
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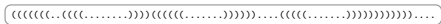
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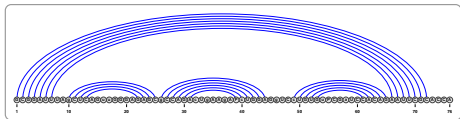




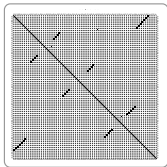
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Hamiltonian-path,  
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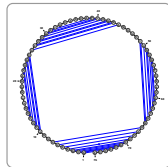
Motzkin words\*



Non-crossing arc-annotated sequences\*



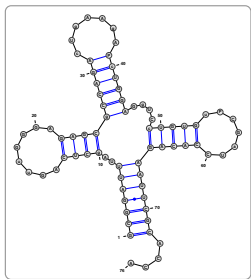
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**Supporting intuitions**  
Different representations  
Common combinatorial structure

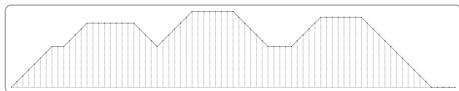
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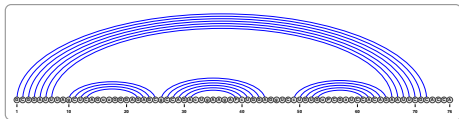
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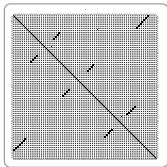
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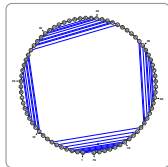
Positive 1D meanders\* over  $S = \{+1, -1, 0\}$



Non-crossing arc-annotated sequences\*



Dot plots  
Adjacency matrices\*



Non-crossing arc diagrams\*

## Supporting intuitions

Different representations  
Common combinatorial structure

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## Theorem

Let  $S(n)$  denote the number of secondary structures of size  $n$  and  $\theta = 1$ . Then  $S(0) = 0$ ,  $S(1) = 1$  and for  $n \geq 1$ ,

$$S(n+1) = S(n) + S(n-1) + \sum_{k=2}^{n-1} S(k-1)S(n-k)$$

and

$$S(n) \geq 2^{n-2}.$$

## Theorem

Let  $S(n, k)$  denote the number of secondary structures of size  $n$  that contain exactly  $k$  basepairs ( $\theta = 1$ ). Set  $S(n, 0) = 1$  for all  $n$  and  $S(n, k) = 0$  for  $k \geq \frac{n}{2}$ . Then for  $n \geq 2$ ,

$$S(n+1, k+1) = S(n, k+1) + \sum_{j=1}^{n-1} \left[ \sum_{i=0}^k S(j-1, i)S(n-j, k-i) \right].$$

## Corollary

$$S(n) = \sum_{k=0}^{\lfloor n/2 \rfloor} S(n, k).$$

All proofs on whiteboard

A sequence  $s \in \mathbb{A}^n$  **realizes** or **is compatible with** a secondary structure  $\mathcal{S}$  of length  $n$  if for any  $(i, j) \in \mathcal{S}$  it holds that  $s_i, s_j \in \mathbb{B}$ .

- ▶ *dependency graph or shape graph*  $G(\mathcal{S}_1, \dots, \mathcal{S}_k)$

### QUESTION:

On what conditions is it possible to find *one* sequence, which is realizing *all* sec.str.  $\mathcal{S}_1, \dots, \mathcal{S}_k$  of the same size?

equivalent to:

What properties have to be fulfilled in  $G(\mathcal{S}_1, \dots, \mathcal{S}_k)$ , s.t. there exists a single sequence, which is realizing all  $\mathcal{S}_1, \dots, \mathcal{S}_k$ ?

$C(\mathcal{S})$  denotes the set of all sequences that realize Secondary Structure  $\mathcal{S}$ .

### Theorem (Intersection Theorem, Reidys et al. 1995)

For any two secondary structures  $\mathcal{S}_1$  and  $\mathcal{S}_2$  of same size holds:  $C(\mathcal{S}_1) \cap C(\mathcal{S}_2) \neq \emptyset$ .

### Theorem (Generalized Intersection Theorem, Flamm et al. 2001)

$\bigcap_{i=1}^k C(\mathcal{S}_i) \neq \emptyset \Leftrightarrow G(\mathcal{S}_1, \dots, \mathcal{S}_k)$  is bipartite.

All proofs on whiteboard

A sequence  $s \in \mathbb{A}^n$  **realizes** or **is compatible with** a secondary structure  $\mathcal{S}$  of length  $n$  if for any  $(i, j) \in \mathcal{S}$  it holds that  $s_i, s_j \in \mathbb{B}$ .

- ▶ *dependency graph or shape graph*  $G(\mathcal{S}_1, \dots, \mathcal{S}_k)$

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**All proofs on whiteboard**

**RNA folding:** hierarchical process in which secondary structure is broadly considered as sufficient approximation assessing the most relevant characteristics of an RNA molecule

$\mathbb{S}(x)$  = set of all second.structures  $\mathcal{S}$  that realize a given RNA sequence  $x$ .

$\mathbb{S}(x)$  is called **folding space** of  $x$ .

Structure prediction means to select the “most-likely” structure from elements of  $\mathbb{S}(x)$ .

“most-likely” = most stable.

Let us start with a simple naive approach, that laid the foundation for many more sophisticated and realistic approaches: **Nussinov Algorithm**



**Nussinov/Jacobson energy model (NJ)****Base-pair maximization** (with a twist):

- ▶ Additive model on **independently contributing** base-pairs;
- ▶ **Canonical base-pairs** only: Watson/Crick (A/U,C/G) and Wobble (G/U)

**Variant:** Weight each pair with #Hydrogen bonds

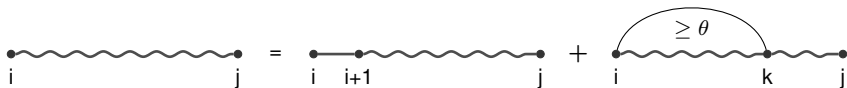
$$w(G\equiv C) = 3 \quad w(A=U) = 2 \quad w(G-U) = 1 \quad w(\text{other}) = 0$$

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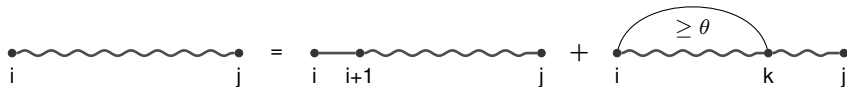
$$w(G \equiv C) = 3 \quad w(A = U) = 2 \quad w(G - U) = 1 \quad w(\text{other}) = 0$$



$$N_{i,t} = 0, \quad \forall t \in [i, i + \theta]$$

$$N_{i,j} = \max \begin{cases} N_{i+1,j} & i \text{ unpaired} \\ \max_{i+\theta < k \leq j} w_{i,k} + N_{i+1,k-1} + N_{k+1,j} & i \text{ paired with } k \end{cases}$$

$$w_{i,k} = w(s_i, s_k) \in \{0, 1, 2, 3\}.$$



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**Correctness.** Goal = Show that MFE over interval  $[i, j]$  is indeed found in  $N_{i,j}$  after completing the computation. Proceed by induction:

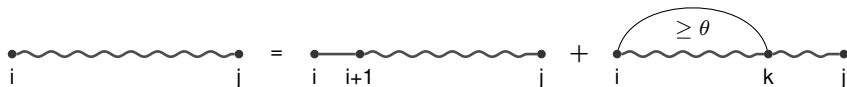
- ▶ Assume that property holds for any  $[i', j']$  such that  $j' - i' < \ell$ .
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Then first position  $i$  in  $\text{MFE}_{i,j}$  is either:

▶ **Unpaired:**  $\text{MFE}_{i,j} = \text{MFE}_{i+1,j} \rightarrow$  free-energy =  $N_{i+1,j}$

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(Indeed, any BP between  $[i+1, k-1]$  and  $[k+1, j]$  would cross  $(i, k)$ )

$\rightarrow$  free-energy =  $w_{i,k} + N_{i+1,k-1} + N_{k+1,j}$



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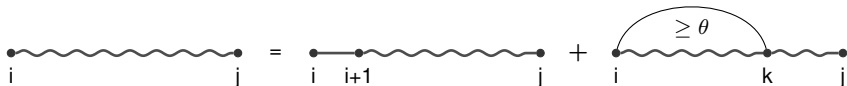
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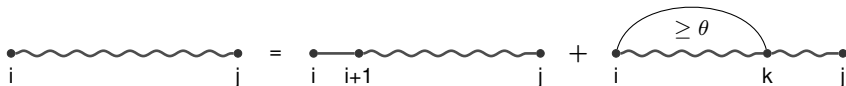
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$$w_{i,k} = w(s_i, s_k) \in \{0, 1, 2, 3\}.$$

$i = 0, j = \text{SIZE} - 1;$

Traceback(matrix  $N, i, j$ )

if  $i < \text{SIZE} - \theta - 1$  then

if  $N_{i,j} = N_{i+1,j}$  then

Traceback( $N, i+1, j$ )

//  $i$  unpaired

else

for( $k = i + \theta + 1, \dots, j$ )

if  $N[i][j] = w_{i,k} + N[i+1][k-1] + N[k+1][j]$  and neither  $i, k$  paired

print "basepair ( $i, k$ )"

Traceback( $N, i+1, k-1$ )

Traceback( $N, k+1, j$ )

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	.	.	.	.	.	.	.	.	.	.	.	.	.	.	.	.	.	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
C							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
C										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A															0	0	0	0
C																0	0	0
G																	0	0
A																		0

Diagram illustrating the Nussinov/Jacobson algorithm's recurrence relation. A sequence of nucleotides from index  $i$  to  $j$  is shown as a wavy line. This is equal to the sum of two terms: 1) a sequence from  $i$  to  $i+1$  followed by a sequence from  $i+1$  to  $j$ , and 2) a sequence from  $i$  to  $k$  followed by a sequence from  $k$  to  $j$ , with a curved arrow above the second term indicating a pairing between  $i$  and  $k$  with energy at least  $\theta$  ( $\geq \theta$ ).



	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	.	.	.	.	.	.	.	.	.	.	.	.	.	.	.	.	.	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
C							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
C										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A															0	0	0	0
C																0	0	0
G																	0	0
A																		0

$i \text{---} j = i \text{---} i+1 \text{---} j + i \text{---} k \text{---} j$

$\geq \theta$

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	.	.	.	.	.	.	.	.	.	.	.	.	.	.	.	.	.	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	2	2	4	5	7	7	8	10	
A						0	0	0	0	2	2	2	5	5	5	8	8	
C							0	0	0	0	0	2	5	5	5	8	8	
U								0	0	0	0	2	3	5	5	6	7	
U									0	0	0	2	3	5	5	5	7	
C										0	0	0	3	3	3	5	5	
U											0	0	0	2	2	2	3	
U												0	0	0	0	1	2	
A													0	0	0	0	0	
G														0	0	0	0	
A															0	0	0	0
C																0	0	0
G																	0	0
A																		0

$$\text{Sequence } (i, \dots, j) = \text{Sequence } (i, \dots, i+1) + \text{Sequence } (i+1, \dots, j) + \text{Sequence } (i, \dots, k) + \text{Sequence } (k, \dots, j)$$

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	.	.	.	.	.	.	.	.	.	.	.	.	.	.	.	.	.	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
C							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
C										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A															0	0	0	0
C																0	0	0
G																	0	0
A																		0

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$\geq \theta$

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A	
	(	.	.	.	.	.	.	.	.	.	.	.	.	.	.	.	)	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14	
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11	
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10	
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10	
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10	
A						0	0	0	0	0	2	2	2	5	5	5	8	8	
C							0	0	0	0	0	0	2	5	5	5	8	8	
U								0	0	0	0	0	2	3	5	5	6	7	
U									0	0	0	0	2	3	5	5	5	7	
C										0	0	0	0	3	3	3	5	5	
U											0	0	0	0	2	2	2	3	
U												0	0	0	0	0	1	2	
A													0	0	0	0	0	0	
G														0	0	0	0	0	
A															0	0	0	0	
C																0	0	0	
G																	0	0	
A																		0	

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(	.	.	.	.	.	.	.	.	.	.	.	.	.	.	.	)	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
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C							0	0	0	0	0	0	2	5	5	5	8	8
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U									0	0	0	0	2	3	5	5	5	7
C										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A															0	0	0	0
C																0	0	0
G																	0	0
A																		0

$i \text{---} j = i \text{---} i+1 \text{---} j + i \text{---} k \text{---} j$

$\geq \theta$

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(	.	.	.	.	.	.	.	.	.	.	.	.	.	.	.	)	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
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C							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
C										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A															0	0	0	0
C																0	0	0
G																	0	0
A																		0

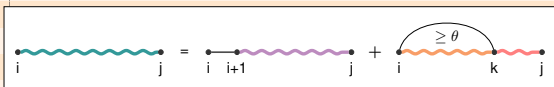
  

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	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(	.	.	.	.	.	.	.	.	.	.	.	.	.	.	.	)	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	2	2	4	5	7	7	8	10	
A						0	0	0	0	2	2	2	5	5	5	8	8	
C							0	0	0	0	0	2	5	5	5	8	8	
U								0	0	0	0	2	3	5	5	6	7	
U									0	0	0	2	3	5	5	5	7	
C										0	0	0	3	3	3	5	5	
U											0	0	0	2	2	3		
U												0	0	0	0	1	2	
A													0	0	0	0	0	
G														0	0	0	0	
A															0	0	0	
C																0	0	0
G																	0	0
A																		0









	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(	.	.	.	.	.	.	.	.	.	.	.	.	.	.	.	)	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
C							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
C										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A															0	0	0	0
C																0	0	0
G																	0	0
A																		0

$$\text{Sequence}(i, j) = \text{Sequence}(i, i+1) + \text{Sequence}(i+1, j) + \text{Sequence}(i, k) + \text{Sequence}(k, j) \quad (\text{with } \theta \text{ constraint})$$

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(	.	.	.	.	.	.	.	.	.	.	.	.	.	.	.	)	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
C							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
C										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A															0	0	0	0
C																0	0	0
G																	0	0
A																		0

$i \text{---} j = i \text{---} i+1 \text{---} j + i \text{---} k \text{---} j$

The diagram shows a sequence of colored wavy lines representing a sequence from index  $i$  to  $j$ . This is decomposed into two parts: a sequence from  $i$  to  $i+1$  (purple) and a sequence from  $i+1$  to  $j$  (pink). Additionally, there is a sequence from  $i$  to  $k$  (orange) with a gap of at least  $\theta$  between  $k$  and  $j$  (red).

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A	
	(	(	.	.	.	.	.	.	.	.	.	.	.	.	.	)	)	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14	
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11	
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10	
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10	
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10	
A						0	0	0	0	0	2	2	2	5	5	5	8	8	
C							0	0	0	0	0	0	2	5	5	5	8	8	
U								0	0	0	0	0	2	3	5	5	6	7	
U									0	0	0	0	2	3	5	5	5	7	
C										0	0	0	0	3	3	3	5	5	
U											0	0	0	0	2	2	2	3	
U												0	0	0	0	0	1	2	
A													0	0	0	0	0	0	
G														0	0	0	0	0	
A															0	0	0	0	
C																0	0	0	
G																	0	0	
A																		0	

$i \text{---} \dots \text{---} j = i \text{---} i+1 \text{---} \dots \text{---} j + i \text{---} \dots \text{---} k \text{---} \dots \text{---} j$ 
  
 where  $k \geq i+1$

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A	
	(	(	.	.	.	.	.	.	.	.	.	.	.	.	.	)	)	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14	
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11	
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10	
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10	
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10	
A						0	0	0	0	0	2	2	2	5	5	5	8	8	
C							0	0	0	0	0	0	2	5	5	5	8	8	
U								0	0	0	0	0	2	3	5	5	6	7	
U									0	0	0	0	2	3	5	5	5	7	
C										0	0	0	0	3	3	3	5	5	
U											0	0	0	0	2	2	2	3	
U												0	0	0	0	0	1	2	
A													0	0	0	0	0	0	
G														0	0	0	0	0	
A															0	0	0	0	
C																0	0	0	
G																	0	0	
A																		0	

Diagram illustrating the decomposition of a sequence from index  $i$  to  $j$ . The sequence is shown as a wavy line. It is equal to the sum of two parts: a sequence from  $i$  to  $i+1$  (a straight line) followed by a sequence from  $i+1$  to  $j$  (a wavy line), plus a sequence from  $i$  to  $k$  (a wavy line) followed by a sequence from  $k$  to  $j$  (a straight line). A curved arrow above the second part indicates a distance of at least  $\theta$  ( $\geq \theta$ ).

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(	(	.	.	.	.	.	.	.	.	.	.	.	.	.	)	)	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	2	2	2	2	4	4	5	7	7	8	10	
U					0	0	0	0	0	2	2	4	5	7	7	8	10	
A						0	0	0	0	2	2	2	5	5	5	8	8	
C							0	0	0	0	0	2	5	5	5	8	8	
U								0	0	0	0	2	3	5	5	6	7	
U									0	0	0	2	3	5	5	5	7	
C										0	0	0	3	3	3	5	5	
U											0	0	0	2	2	2	3	
U												0	0	0	0	1	2	
A													0	0	0	0	0	
G														0	0	0	0	
A															0	0	0	
C																0	0	0
G																	0	0
A																		0

Diagram illustrating a recurrence relation for sequence alignment. A green wavy line from index  $i$  to  $j$  is equal to a purple wavy line from  $i$  to  $i+1$  plus a red wavy line from  $i$  to  $k$  to  $j$ , where the red line is under a curved arrow labeled  $\geq \theta$ .

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(	(	(	.	.	.	)	.	.	.	.	.	.	.	.	)	)	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	0	2	2	2	2	4	4	5	7	7	8	10
U					0	0	0	0	0	0	2	2	4	5	7	7	8	10
A						0	0	0	0	0	2	2	2	5	5	5	8	8
C							0	0	0	0	0	0	2	5	5	5	8	8
U								0	0	0	0	0	2	3	5	5	6	7
U									0	0	0	0	2	3	5	5	5	7
C										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A															0	0	0	0
C																0	0	0
G																	0	0
A																		0

$$\text{Sequence}(i, j) = \text{Sequence}(i, i+1) + \text{Sequence}(i+1, j) + \sum_{k=i+\theta}^j \text{Sequence}(i, k) + \text{Sequence}(k, j)$$



	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(	(	(	.	.	.	)	.	.	.	.	.	.	.	.	)	)	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	2	2	2	2	4	4	5	7	7	8	10	
U					0	0	0	0	0	2	2	4	5	7	7	8	10	
A						0	0	0	0	2	2	2	5	5	5	8	8	
C							0	0	0	0	0	2	5	5	5	8	8	
U								0	0	0	0	2	3	5	5	6	7	
U									0	0	0	2	3	5	5	5	7	
C										0	0	0	3	3	3	5	5	
U											0	0	0	2	2	2	3	
U												0	0	0	0	1	2	
A													0	0	0	0	0	
G														0	0	0	0	
A															0	0	0	
C																0	0	0
G																	0	0
A																		0

$$\text{Sequence}(i, j) = \text{Sequence}(i, i+1) + \text{Sequence}(i+1, j) + \sum_{i < k < j, k-i \geq \theta} \text{Sequence}(i, k) + \text{Sequence}(k, j)$$

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A	
	(	(	(	.	.	.	)	.	.	.	.	.	.	.	.	)	)	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14	
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11	
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10	
A				0	0	0	2	2	2	2	4	4	5	7	7	8	10		
U					0	0	0	0	0	2	2	4	5	7	7	8	10		
A						0	0	0	0	2	2	2	5	5	5	8	8		
C							0	0	0	0	0	2	5	5	5	8	8		
U								0	0	0	0	2	3	5	5	6	7		
U									0	0	0	2	3	5	5	5	7		
C										0	0	0	3	3	3	5	5		
U											0	0	0	2	2	2	3		
U												0	0	0	0	1	2		
A													0	0	0	0	0		
G														0	0	0	0		
A															0	0	0		
C																0	0		
G																	0		
A																		0	

$i \text{---} j = i \text{---} i+1 \text{---} j + i \text{---} k \text{---} j$

$\geq \theta$



	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A	
	(	(	(	.	.	.	)	.	.	.	.	.	.	.	.	)	)	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14	
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11	
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10	
A				0	0	0	2	2	2	2	4	4	5	7	7	8	10		
U					0	0	0	0	0	2	2	4	5	7	7	8	10		
A						0	0	0	0	2	2	2	5	5	5	8	8		
C							0	0	0	0	0	2	5	5	5	8	8		
U								0	0	0	0	2	3	5	5	6	7		
U									0	0	0	2	3	5	5	5	7		
C										0	0	0	3	3	3	5	5		
U											0	0	0	2	2	2	3		
U												0	0	0	0	1	2		
A													0	0	0	0	0		
G														0	0	0	0		
A															0	0	0		
C																0	0		
G																	0		
A																		0	

$i \text{---} j = i \text{---} i+1 \text{---} j + i \text{---} k \text{---} j$

$\geq \theta$



	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A	
	(	(	(	.	.	.	)	.	.	.	.	.	.	.	.	)	)	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14	
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11	
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10	
A				0	0	0	2	2	2	2	4	4	5	7	7	8	10		
U					0	0	0	0	0	2	2	4	5	7	7	8	10		
A						0	0	0	0	2	2	2	5	5	5	8	8		
C							0	0	0	0	0	2	5	5	5	8	8		
U								0	0	0	0	2	3	5	5	6	7		
U									0	0	0	2	3	5	5	7			
C										0	0	0	0	3	3	5	5		
U											0	0	0	2	2	3			
U												0	0	0	0	1	2		
A													0	0	0	0			
G														0	0	0			
A															0	0			
C																0			
G																	0		
A																		0	

$$\text{Green}(i, j) = \text{Purple}(i, i+1) + \text{Red}(i, k, j)$$

The diagram shows a sequence from index  $i$  to  $j$ . The first part is a green wavy line. This is equal to the sum of a purple wavy line from  $i$  to  $i+1$  and a red wavy line from  $i$  to  $k$  to  $j$ . A curved arrow above the red line is labeled  $\geq \theta$ .

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A	
	(	(	(	.	.	.	)	.	(	.	.	.	.	.	)	)	)	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14	
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11	
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10	
A				0	0	0	2	2	2	2	4	4	5	7	7	8	10		
U					0	0	0	0	0	2	2	4	5	7	7	8	10		
A						0	0	0	0	2	2	2	5	5	5	8	8		
C							0	0	0	0	0	2	5	5	5	8	8		
U								0	0	0	0	2	3	5	5	6	7		
U									0	0	0	2	3	5	5	5	7		
C										0	0	0	3	3	3	5	5		
U											0	0	0	2	2	2	3		
U												0	0	0	0	1	2		
A													0	0	0	0	0		
G														0	0	0	0		
A															0	0	0		
C																0	0	0	
G																	0	0	
A																		0	

$$\text{Sequence}(i, j) = \text{Sequence}(i, i+1) + \text{Sequence}(i+1, j) + \sum_{k=i+\theta}^j \text{Sequence}(i, k) + \text{Sequence}(k, j)$$

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A
	(	(	(	.	.	.	)	.	(	.	.	.	.	.	)	)	)	.
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10
A				0	0	0	2	2	2	2	4	4	5	7	7	8	10	
U					0	0	0	0	0	2	2	4	5	7	7	8	10	
A						0	0	0	0	2	2	2	5	5	5	8	8	
C							0	0	0	0	0	2	5	5	5	8	8	
U								0	0	0	0	2	3	5	5	6	7	
U									0	0	0	2	3	5	5	5	7	
C										0	0	0	0	3	3	3	5	5
U											0	0	0	0	2	2	2	3
U												0	0	0	0	0	1	2
A													0	0	0	0	0	0
G														0	0	0	0	0
A															0	0	0	0
C																0	0	0
G																	0	0
A																		0

$$\text{wavy}(i, j) = \text{wavy}(i, i+1) + \text{wavy}(i+1, j) + \sum_{k=i+\theta}^j \text{wavy}(i, k) + \text{wavy}(k, j)$$



	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A	
	(	(	(	.	.	.	)	.	(	.	.	.	.	.	)	)	)	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14	
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11	
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10	
A				0	0	0	2	2	2	2	4	4	5	7	7	8	10		
U					0	0	0	0	0	2	2	4	5	7	7	8	10		
A						0	0	0	0	2	2	2	5	5	5	8	8		
C							0	0	0	0	0	2	5	5	5	8	8		
U								0	0	0	0	2	3	5	5	6	7		
U									0	0	0	2	3	5	5	5	7		
C										0	0	0	3	3	3	5	5		
U											0	0	0	2	2	2	3		
U												0	0	0	0	1	2		
A													0	0	0	0	0		
G														0	0	0	0		
A															0	0	0		
C																0	0	0	
G																	0	0	
A																		0	

$$\text{wavy}(i, j) = \text{wavy}(i, i+1) + \text{wavy}(i, k) + \text{wavy}(k, j)$$

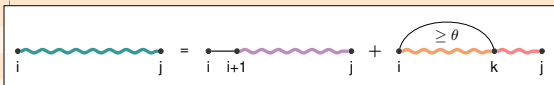
The diagram shows a wavy line from index  $i$  to  $j$  on the left. This is equal to the sum of two terms on the right: a wavy line from  $i$  to  $i+1$ , and a wavy line from  $i$  to  $k$  with a loop above it labeled  $\geq \theta$ , plus a wavy line from  $k$  to  $j$ .

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A	
	(	(	(	.	.	.	)	.	(	(	.	.	.	)	)	)	)	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14	
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11	
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10	
A				0	0	0	2	2	2	2	4	4	5	7	7	8	10		
U					0	0	0	0	0	2	2	4	5	7	7	8	10		
A						0	0	0	0	2	2	2	5	5	5	8	8		
C							0	0	0	0	0	2	5	5	5	8	8		
U								0	0	0	0	2	3	5	5	6	7		
U									0	0	0	2	3	5	5	5	7		
C										0	0	0	3	3	3	5	5		
U											0	0	0	2	2	2	3		
U												0	0	0	0	1	2		
A													0	0	0	0	0		
G														0	0	0	0		
A															0	0	0		
C																0	0	0	
G																	0	0	
A																		0	

$$\text{Sequence}(i, j) = \text{Sequence}(i, i+1) + \text{Sequence}(i+1, j) + \sum_{k=i+1}^j \text{Sequence}(i, k) + \text{Sequence}(k, j)$$

	C	G	G	A	U	A	C	U	U	C	U	U	A	G	A	C	G	A	
	(	(	(	.	.	.	)	.	(	(	.	.	.	)	)	)	)	.	
C	0	0	0	0	0	0	3	4	4	6	6	6	6	9	9	11	14	14	
G		0	0	0	0	0	3	4	4	6	6	6	6	7	9	11	11	11	
G			0	0	0	0	3	3	3	5	5	5	5	6	8	10	10	10	
A				0	0	0	2	2	2	2	4	4	5	7	7	8	10		
U					0	0	0	0	0	2	2	4	5	7	7	8	10		
A						0	0	0	0	2	2	2	5	5	5	8	8		
C							0	0	0	0	0	2	5	5	5	8	8		
U								0	0	0	0	2	3	5	5	6	7		
U									0	0	0	2	3	5	5	5	7		
C										0	0	0	3	3	3	5	5		
U											0	0	0	2	2	2	3		
U												0	0	0	0	1	2		
A													0	0	0	0	0		
G														0	0	0	0		
A															0	0	0		
C																0	0	0	
G																	0	0	
A																		0	



Maximizing the nr of bp does not lead to biological meaningful structures:

Stacking of bp not considered:

( ( ( . ) ) )    **stable**  
 ( ) ( ) . ( )    **instable**

Size of intern loop not considered:

instable



stable

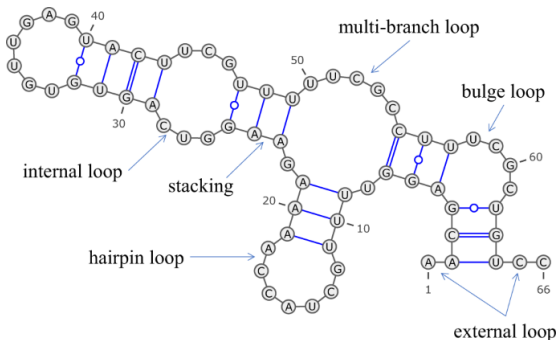


instable



Nevertheless, although Nussinov-Alg is too simple to be accurate, it is stepping-stone for later algorithms

Define energy model for RNA that takes into account local energy contributions from loop and stacking regions.



- ▶ More realistic: thermodynamics and statistical mechanics.
- ▶ Stability of an RNA sec.str. coincides with thermodynamic stability
- ▶ Quantified as the amount of free energy released/used by forming bp.

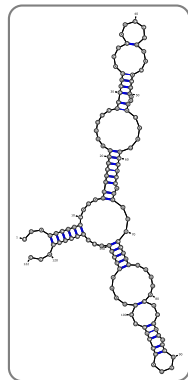
Based on **unambiguous** decomposition of  $2^{\text{ary}}$  structure into **loops**:

- ▶ Internal loops
- ▶ Bulges
- ▶ Terminal (hairpin) loops
- ▶ Multi loops
- ▶ Stackings

Free-energy  $\Delta G$  of a loop depend on bases, asymmetry, dangles ...

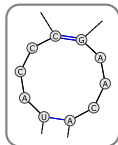
Experimentally determined  
+ Interpolated for larger loops.

Improved results by taking stacking into account.



Based on **unambiguous** decomposition of  $2^{\text{ary}}$  structure into **loops**:

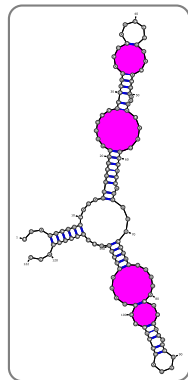
- ▶ Internal loops
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Free-energy  $\Delta G$  of a loop depend on bases, asymmetry, dangles ...

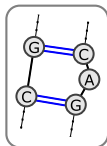
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Based on **unambiguous** decomposition of  $2^{\text{ary}}$  structure into **loops**:

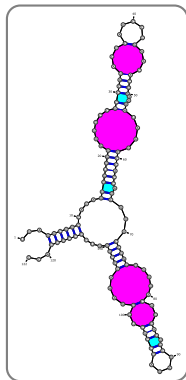
- ▶ Internal loops
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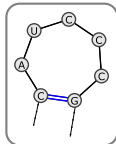
Improved results by taking stacking into account.





Based on **unambiguous** decomposition of  $2^{\text{ary}}$  structure into **loops**:

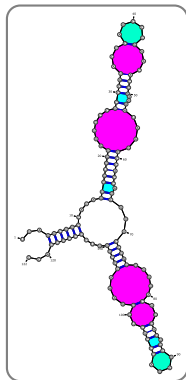
- ▶ Internal loops
- ▶ Bulges
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Free-energy  $\Delta G$  of a loop depend on bases, asymmetry, dangles ...

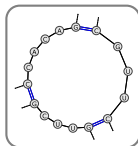
Experimentally determined  
+ Interpolated for larger loops.

Improved results by taking stacking into account.



Based on **unambiguous** decomposition of  $2^{\text{ary}}$  structure into **loops**:

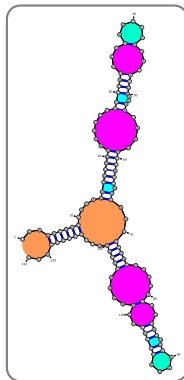
- ▶ Internal loops
- ▶ Bulges
- ▶ Terminal (hairpin) loops
- ▶ Multi loops
- ▶ Stackings



Free-energy  $\Delta G$  of a loop depend on bases, asymmetry, dangles ...

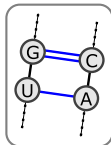
Experimentally determined  
+ Interpolated for larger loops.

Improved results by taking stacking into account.



Based on **unambiguous** decomposition of  $2^{\text{ary}}$  structure into **loops**:

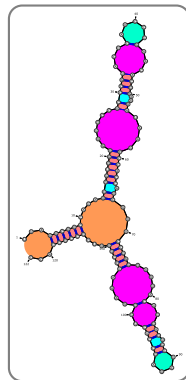
- ▶ Internal loops
- ▶ Bulges
- ▶ Terminal (hairpin) loops
- ▶ Multi loops
- ▶ Stackings



Free-energy  $\Delta G$  of a loop depend on bases, asymmetry, dangles ...

Experimentally determined  
+ Interpolated for larger loops.

Improved results by taking stacking into account.



The Turner rules are a set of experimentally determined parameters which allow us to predict the stability of RNA secondary structures.

## Turner Energy Rules

		TOP					
		AU	CG	GC	UA	GU	UG
B O T T O M	AU	-0.9	-1.8	-2.3	-1.1	-0.5	-0.7
	CG	-2.1	-2.9	-3.4	-2.3	-1.5	-1.5
	GC	-1.7	-2	-2.9	-1.8	-1.3	-1.5
	UA	-0.9	-1.7	-2.1	-0.9	-0.7	-0.5
	GU	-0.9	-1.7	-2.1	-0.9	-0.5	-0.5
	UG	-0.9	-1.7	-2.1	-0.9	0.6	-0.5

Bases in Loop	Internal Loop	Bulge Loop	Hairpin Loop
1	0	3.3	0
2	0.8	5.2	0
3	1.3	6	7.4
4	1.7	6.7	5.9
5	2.1	7.4	4.4
6	2.5	8.2	4.3
7	2.6	9.1	4.1
8	2.8	10	4.1

RESET

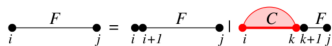
PRACTICE

PRINT EXAM

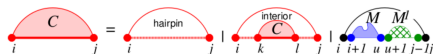
```

      G
     C   C
    A   C
   AU
  GU
 G   G
A   A
  GU
 C |
  UA
  GU
U |
  CG
 C |
  AU
U   U
A   C
  GU
5'C A3' RNA Free Energy:

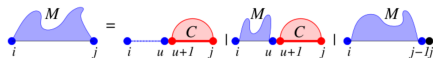
```



$$F_{i,j} = \min \left\{ \begin{array}{l} F_{i+1,j}, \\ \min_{i < k \leq j} C_{i,k} + F_{k+1,j} \end{array} \right\}$$



$$C_{i,j} = \min \left\{ \begin{array}{l} \mathcal{H}(i,j), \\ \min_{i < k < l < j} C_{k,l} + \mathcal{I}(i,j;k,l), \\ \min_{i < u < j} M_{i+1,u} + M_{u+1,j-1}^1 + a \end{array} \right\}$$



$$M_{i,j} = \min \left\{ \begin{array}{l} \min_{i < u < j} (u - i + 1)c + C_{u+1,j} + b, \\ \min_{i < u < j} M_{i,u} + C_{u+1,j} + b, \\ M_{i,j-1} + c \end{array} \right\}$$



$$M^1_{i,j} = \min \left\{ \begin{array}{l} M^1_{i,j-1} + c, \\ C_{i,j} + b \end{array} \right\}$$

First proposed by Zuker et al.

**init:**  $F_{i,j} = 0$ ,  $C_{i,j} = M_{i,j} = M^1_{i,j} = \infty$ .

- ▶  $F_{1,n}$  stores the energy value of the thermodynamically most stable structure, its Minimum Free Energy (MFE).
- ▶ traceback structure

- ▶  $F_{i,j}$  : free energy of the opt. sub-struct. on the sub-seq.  $s_i \dots s_j$ .
- ▶  $C_{i,j}$  : free energy of the opt. sub-struct. on the sub-seq.  $s_i \dots s_j$  given that  $i$  and  $j$  form a base pair.
- ▶  $M_{i,j}$  : free energy of the opt. sub-struct. on the sub-seq.  $s_i \dots s_j$  given that  $s_i \dots s_j$  is part of a multi-loop and has at least one "component".
- ▶  $M_{i,j}^1$  : free energy of the opt. sub-struct. on the sub-seq.  $s_i \dots s_j$  given that  $s_i \dots s_j$  is part of a multi-loop and has exactly one component which has the closing pair  $(i, h)$  for some  $h$  satisfying  $i < h \leq j$ .

```
http://rna.tbi.univie.ac.at/cgi-bin/RNAfold.cgi
```

```
RNAfold < trna.fa  
>AF041468  
GGGGGUUAGCUCAGUUGGUAGAGCGCUGCCUUUGCACGGCAGAUGUCAGGGGUUCGAGUCCCCUACCUCCA  
(((((((..((((.....))))).((((.....))))).(((.....)))))))).  
-31.10 kcal/mol
```





